

# Cone avoidance and randomness preservation

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## Basis theorems.

A basis theorem is a theorem of the form:

For any nonempty effectively closed set in Euclidean space, at least one member of the set is “close to being computable”.

Some well known basis theorems are:

- the Low Basis Theorem,
- the R.E. Basis Theorem,
- the Hyperimmune-Free Basis Theorem,
- the Cone Avoidance Basis Theorem,
- the Randomness Preservation Basis Thm.

Less well known is a basis theorem of Higuchi/Hudelson/Simpson/Yokoyama on preservation of partial randomness.

Basis theorems are important for applications in the foundations of mathematics: models of arithmetic, Scott sets,  $\omega$ -models of  $WKL_0$  and  $WWKL_0$ , reverse mathematics, etc.

We discuss the possibilities for combining these basis theorems.

## Three basis theorems.

Let  $\leq_T$  denote Turing reducibility.

Let  $'$  denote the Turing jump operator.

### The Low Basis Theorem:

For any nonempty effectively closed set  $Q$ , there exists  $Z \in Q$  such that  $Z' \leq_T 0'$ .

### The R.E. Basis Theorem:

For any nonempty effectively closed set  $Q$ , there exists  $Z \in Q$  such that  $Z$  is of recursively enumerable Turing degree.

We say that  $Z$  is *hyperimmune-free* if  $(\forall \text{ functions } f \leq_T Z) (\exists \text{ recursive function } g) \forall n (f(n) < g(n))$ .

### The Hyperimmune-Free Basis Theorem:

For any nonempty effectively closed set  $Q$ ,  $(\exists Z \in Q) (Z \text{ is hyperimmune-free})$ .

These three basis theorems are due to Jockusch/Soare 1972.

## Can we combine these basis theorems?

No. The Jockusch/Soare basis theorems are known to be “pairwise incompatible.”

1. The Arslanov Completeness Criterion provides a nonempty effectively closed  $Q$  such that for all r.e. sets  $A$ , if  $(\exists Z \in Q) (Z \leq_T A)$  then  $0' \leq_T A$ .

Therefore, the Low Basis Theorem and the R.E. Basis Theorem cannot be combined into one basis theorem.

2. It is known that for hyperimmune-free  $Z$  one cannot have  $0 <_T Z \leq_T 0'$ .

Therefore, the Hyperimmune-Free Basis Theorem cannot be combined with the Low Basis Theorem or with the R.E. Basis Theorem.

## Two more basis theorems.

### The Cone Avoidance Basis Theorem:

For any nonempty effectively closed set  $Q$ ,  
if  $A \not\leq_T 0$  then  $(\exists Z \in Q) (A \not\leq_T Z)$ .

More generally,

if  $\forall i (A_i \not\leq_T 0)$  then  $(\exists Z \in Q) \forall i (A_i \not\leq_T Z)$ .

Gandy/Kreisel/Tait, 1960.

Let  $\text{MLR} = \{X \mid X \text{ is Martin-Löf random}\}$ .

Let  $\text{MLR}^Z = \{X \mid X \text{ is Martin-Löf random relative to } Z\}$ .

### The Randomness Preservation Basis Theorem:

For any nonempty effectively closed set  $Q$ ,  
if  $X \in \text{MLR}$  then  $(\exists Z \in Q) (X \in \text{MLR}^Z)$ .

Reimann/Slaman, 2005,

Downey/Hirschfeldt/Miller/Nies, 2005,

Simpson/Yokoyama, 2011.

## More combinations of basis theorems?

It is known that Cone Avoidance can be combined with the Low Basis Theorem, or with the Hyperimmune-free Basis Theorem, but not with the R.E. Basis Theorem. (See for instance Downey/Hirschfeldt §2.19.3.)

Also, Randomness Preservation cannot be combined with the Low or the R.E. or the Hyperimmune-Free Basis Theorem.

Specifically, let  $\Omega \in \text{MLR}$  be such that  $\Omega \equiv_{\top} 0'$ . It is known that such reals exist (Chaitin, Kučera/Gács). We then have:

1. Any  $Z \leq_{\top} 0'$  such that  $\Omega \in \text{MLR}^Z$  is K-trivial, hence not PA-complete. (See Chapter 11 of Downey/Hirschfeldt 2010 or Chapter 5 of Nies 2009.)

2. Any hyperimmune-free  $Z$  such that  $\Omega \in \text{MLR}^Z$  is recursive. (See Theorem 8.1.18 of Nies 2009.)

## Combining basis theorems.

	Low	R.E.	H.I.F.	C.A.	R.P.
Low	1	0	0	1	0
R.E.	0	1	0	0	0
H.I.Free	0	0	1	1	0
Cone Av.	1	0	1	1	???
Rand. Pres.	0	0	0	???	1

**Remaining question:** Can Cone Avoidance be combined with Randomness Preservation?

The answer to this question involves LR-reducibility.

Define  $A \leq_{LR} B \iff \text{MLR}^B \subseteq \text{MLR}^A$ . Clearly  $A \leq_T B$  implies  $A \leq_{LR} B$ , and it is known that  $A \leq_{LR} 0$  implies  $A' \leq_T 0'$ . A major theorem of Nies is that  $A \leq_{LR} 0 \iff A$  is K-trivial. See Nies 2009 or Downey/Hirschfeldt 2010.

A theorem which combines Cone Avoidance and Randomness Preservation:

**Theorem 1** (Simpson/Stephan, 2013).

For any nonempty effectively closed set  $Q$ , if  $X \in \text{MLR}$  and  $\forall i (A_i \not\leq_{\text{LR}} 0$  or  $A_i \not\leq_{\text{T}} X)$ , then  $(\exists Z \in Q) (X \in \text{MLR}^Z$  and  $\forall i (A_i \not\leq_{\text{T}} Z))$ .

On the other hand, let  $\Omega \in \text{MLR}$  be such that  $\Omega \equiv_{\text{T}} 0'$ . It is well known that such reals exist (Chaitin, Kučera/Gács).

**Theorem 2** (Simpson/Stephan, 2013).

$\exists$  nonempty effectively closed set  $Q$  such that  $(\forall A \leq_{\text{LR}} 0) (\forall Z \in Q) (\Omega \in \text{MLR}^Z \Rightarrow A \leq_{\text{T}} Z)$ .

## Summary of Theorems 1 and 2:

Cone Avoidance is “almost compatible” with Randomness Preservation.

The only obstacle to full compatibility is the existence of non-computable K-trivial cones, i.e.,  $A \leq_{\text{LR}} 0$  and  $A \not\leq_{\text{T}} 0$ .



## Proofs of Theorems 1 and 2.

To prove Theorem 1, we use the Cone Avoidance Basis Theorem, relativized to  $X$ .

To prove Theorem 2, we use  $K =$  prefix-free Kolmogorov complexity.

(1) If  $\Omega \in \text{MLR}^Z$  then  $|K(n) - K^Z(n)| \leq O(1)$  for infinitely many  $n$ . (Miller 2010.)

(2) If  $\Omega \in \text{MLR}^Z$  and  $Z$  is PA-complete, then there exist a  $Z$ -recursive function  $F$  and an infinite  $Z$ -recursive set  $A$  such that  $|K(n) - F(n)| \leq O(1)$  for all  $n \in A$ .

(3) Let  $C =$  plain Kolmogorov complexity. Chaitin 1976 proved: every  $C$ -trivial real is computable. Using  $F$  and  $A$  as in (2), we similarly prove: every  $K$ -trivial real is  $\leq_T Z$ .

For details, see Simpson/Stephan 2013.

## An application.

Recall that  $WKL_0$  is a subsystem of  $Z_2$  which is good for the reverse mathematics of compactness (Heine-Borel, Arzela-Ascoli, Hahn-Banach, fixed points, prime ideals, ...).

And,  $WWKL_0$  is a subsystem of  $WKL_0$  which is good for the reverse mathematics of measure theory (countable additivity, Monotone and Dominated Convergence theorems, Vitali Covering Lemma, ...).

Let  $M$  be a countable  $\omega$ -model of  $WWKL_0$ .

By Simpson/Yokoyama 2011, we get a countable  $\omega$ -model  $M_1 \supseteq M$  of  $WKL_0$  such that  $C \cap M \neq \emptyset$  for every  $M_1$ -coded closed set  $C$  of positive measure.

Call this a good extension of  $M$ .

As an application of Theorem 1, we get two good extensions  $M_1, M_2 \supseteq M$  such that  $M = M_1 \cap M_2$ .

## Partial randomness.

Let  $f : \{0, 1\}^* \rightarrow [0, \infty)$  be computable.

For  $S \subseteq \{0, 1\}^*$  let  $\text{wt}_f(S) = \sum_{\sigma \in S} 2^{-f(\sigma)}$ ,

$\text{pwt}_f(S) = \sup\{\text{wt}_f(P) \mid P \subseteq S \text{ prefix-free}\}$ ,

and  $\llbracket S \rrbracket = \{X \in \{0, 1\}^{\mathbb{N}} \mid \exists n (X \upharpoonright n \in S)\}$ .

We say that  $X$  is strongly  $f$ -random if

$X \notin \bigcap_{i=0}^{\infty} \llbracket S_i \rrbracket$  for all uniformly r.e.  $S_i \subseteq \{0, 1\}^*$

such that  $\forall i (\text{pwt}_f(S_i) \leq 2^{-i})$ .

Martin-Löf randomness is the special case  $f(\sigma) = |\sigma|$ . In this case  $\text{pwt}_f(S) = \lambda(\llbracket S \rrbracket)$  where  $\lambda$  is the fair coin measure on  $\{0, 1\}^{\mathbb{N}}$ .

**Theorem** (Partial Randomness Preservation, Higuchi/Hudelson/Simpson/Yokoyama 2011).

For any nonempty effectively closed set  $Q$ , if  $X$  is strongly  $f$ -random then  $(\exists Z \in Q)$  ( $X$  is strongly  $f$ -random relative to  $Z$ ).

More generally, if  $\forall i (X_i \text{ is strongly } f_i\text{-random})$  then  $(\exists Z \in Q) \forall i (X_i \text{ is strongly } f_i\text{-random relative to } Z)$ .

**Problem:** To what extent can we combine the *Partial* Randomness Preservation Basis Theorem with cone avoidance?

**Theorem 3** (implicit in H/H/S/Y 2011).

For any nonempty effectively closed set  $Q$ , if  $\forall i (A_i \not\leq_{LR} 0$  and  $X_i$  is strongly  $f_i$ -random), then  $(\exists Z \in Q) \forall i (A_i \not\leq_{LR} Z$  and  $X_i$  is strongly  $f_i$ -random relative to  $Z$ ).

On the other hand, because of Theorem 2, we cannot *always* replace  $\leq_{LR}$  by  $\leq_T$ .

Can we *sometimes* replace  $\leq_{LR}$  by  $\leq_T$ ?

**A typical question:**

Define  $X$  to be strongly half-random  $\iff$   $X$  is strongly  $f$ -random where  $f(\sigma) = |\sigma|/2$ .

Let  $Q$  be nonempty and effectively closed.

If  $A \not\leq_T 0$  and  $X$  is strongly half-random, does there exist  $Z \in Q$  such that  $A \not\leq_T Z$  and  $X$  is strongly half-random relative to  $Z$ ?

## Recent literature.

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Stephen G. Simpson and Frank Stephan, Cone avoidance and randomness preservation, 22 pages, 2013, Annals of Pure and Applied Logic, conditionally accepted for publication.

## More on partial randomness.

Under mild hypotheses on  $f$ , Hudelson proved the existence of a strong  $f$ -random which does not compute any (strong)  $g$ -random, if  $g$  grows significantly faster than  $f$ .

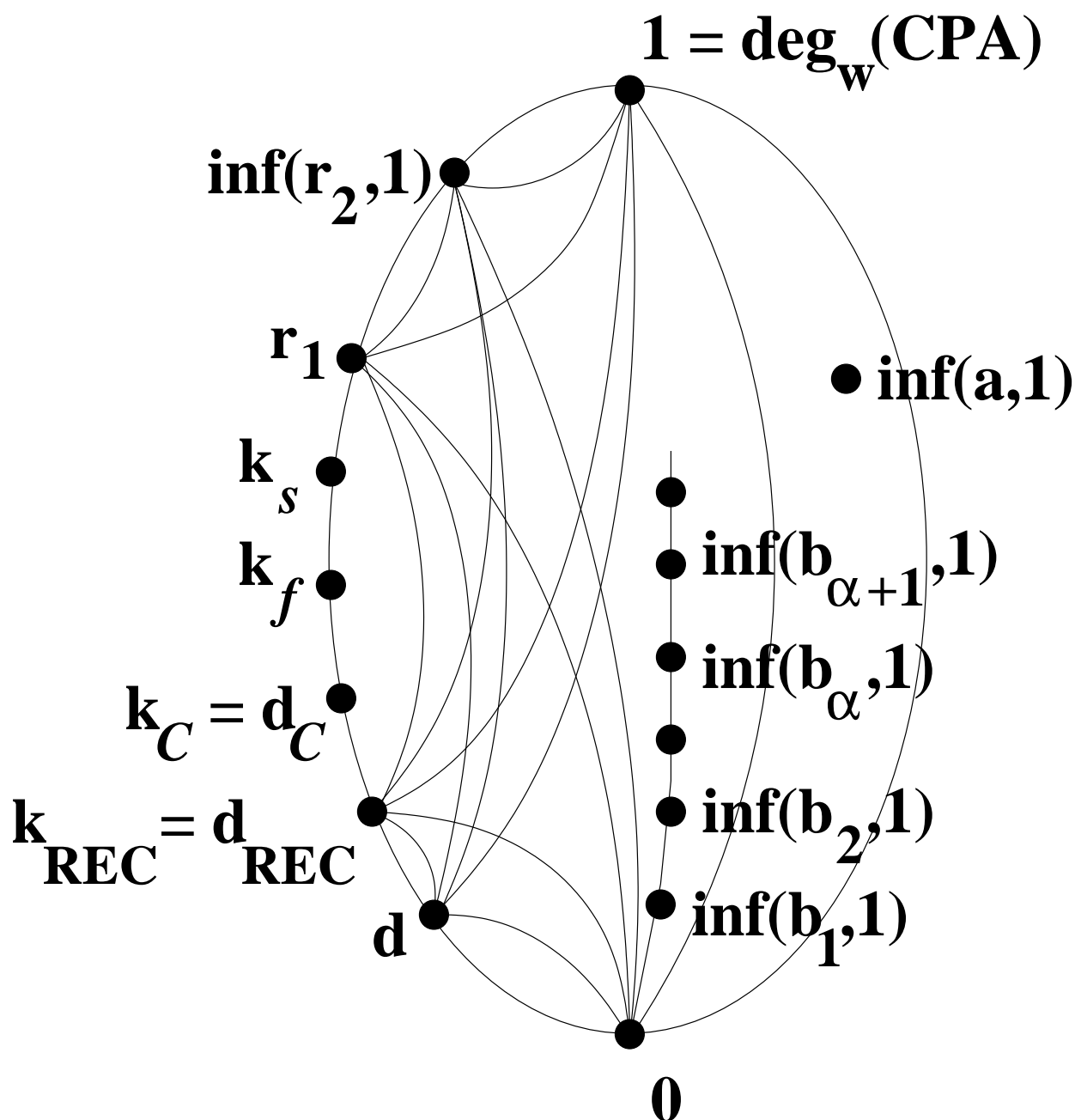
**Theorem** (Hudelson). Let  $f$  and  $g$  be computable, convex, unbounded, and length-invariant such that  $f(\sigma) + 2 \log_2 f(\sigma) \leq g(\sigma)$  for all  $\sigma$ . Then  $\exists X$  ( $X$  is strongly  $f$ -random and  $(\forall Y \leq_T X)$  ( $Y$  is not  $g$ -random)).

This generalizes results of Miller and Greenberg/Miller.

**Example.** We can find an  $X$  such that  $K(X \upharpoonright n) \geq^+ \sqrt{n}$  and there is no  $Y \leq_T X$  such that  $K(Y \upharpoonright n) \geq^+ \sqrt{n} + \log_2 n$ .

**Reference.** W. M. Phillip Hudelson, Mass problems and initial segment complexity, *Journal of Symbolic Logic*, 79, 2014, 20–44.

**Note.** Hudelson's theorem provides many natural examples of Muchnik degrees of mass problems associated with nonempty  $\Pi_1^0$  subsets of  $\{0, 1\}^{\mathbb{N}}$ .



## Partial randomness and $\mu$ -randomness.

Several authors (Levin/Gács, Day/Miller, Reimann/Slaman, Day/Reimann, . . . ) have defined what it means for  $X \in \{0, 1\}^{\mathbb{N}}$  to be  $\mu$ -random where  $\mu$  is an arbitrary Borel probability measure on  $\{0, 1\}^{\mathbb{N}}$ .

If  $\mu = \lambda =$  the fair coin measure on  $\{0, 1\}^{\mathbb{N}}$ , then  $\mu$ -randomness = Martin-Löf randomness.

In general,  $\mu$  need not be computable.

From now on, let  $f : \{0, 1\}^* \rightarrow [0, \infty)$  be computable and convex, i.e.,  $\text{wt}_f(\sigma) \leq \text{wt}_f(\sigma \hat{\ } \langle 0 \rangle) + \text{wt}_f(\sigma \hat{\ } \langle 1 \rangle)$  for all  $\sigma$ . This is a mild assumption.

We can characterize strong  $f$ -randomness in terms of  $\mu$ -randomness:

### Effective Capacitability Theorem

(Reimann, Kjos-Hanssen, Simpson/Stephan).

$X$  is strongly  $f$ -random  $\iff$

$\exists \mu (X \text{ is } \mu\text{-random} \wedge \exists c \forall \sigma (\mu(\llbracket \sigma \rrbracket) \leq 2^{c-f(\sigma)}))$ .



A product theorem for  $\mu \times \nu$ -randomness:

**Theorem.**  $X \oplus Y$  is  $\mu \times \nu$ -random  $\iff X$  is  $\mu$ -random and  $Y$  is  $\nu$ -random relative to  $X, \mu$ .

Combine with Effective Capacitability:

**Theorem 4** (Simpson/Stephan 2013).

If  $X$  is strongly  $f$ -random, and  
if  $Y$  is Martin-Löf random relative to  $X$ ,  
then  $X$  is strongly  $f$ -random relative to  $Y$ .

Theorem 4 resembles an older result:

**Theorem** (H/H/S/Y 2011).

If  $X$  is strongly  $f$ -random and  $\leq_T Y$   
where  $Y$  is Martin-Löf random relative to  $Z$ ,  
then  $X$  is strongly  $f$ -random relative to  $Z$ .

However, we do not know how to deduce  
Theorem 4 from H/H/S/Y or vice versa.

Jason Rute has used Effective Capacitability  
to prove an Ample Excess Lemma for strong  
 $f$ -randomness:

**Theorem** (Rute). If  $X$  is strongly  $f$ -random,

then  $\sum_{n=0}^{\infty} 2^{f(X \upharpoonright n) - K(X \upharpoonright n)} < \infty$ .

## Complexity and autocomplexity.

**Definition** (Kjos-Hanssen/Merkle/Stephan).  
 $X \in \{0, 1\}^{\mathbb{N}}$  is complex (autocomplex)  
if there exists an unbounded  $h : \mathbb{N} \rightarrow \mathbb{N}$   
such that  $K(X \upharpoonright n) \geq h(n)$  for all  $n$ , and  
 $h$  is computable (computable from  $X$ ).

**Theorem** (K-H/M/S 2006).  
 $X$  is complex (autocomplex)  $\iff$   
there exists a DNR function  $\leq_{\text{tt}} X$  ( $\leq_{\text{T}} X$ ).

**Theorem** (H/H/S/Y 2011).  
 $X$  is autocomplex (complex)  $\iff$   
 $X$  is strongly  $f$ -random for some  
computable (computable length-invariant)  
 $f$  such that  $\{f(X \upharpoonright n) \mid n \in \mathbb{N}\}$  is unbounded.

**Theorem** (Reimann/Slaman, Simpson/Stephan).

1.  $X$  is autocomplex relative to some oracle  
 $\iff \exists \mu (X \text{ is } \mu\text{-random} \wedge \mu(\{X\}) = 0),$   
 $\iff X$  is non-computable.
2.  $X$  is complex relative to some oracle  
 $\iff \exists \mu (X \text{ is } \mu\text{-random} \wedge \forall Y (\mu(\{Y\}) = 0)).$

## Literature.

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**Thank you for your attention!**