

Weak Lowness Notions for Kolmogorov Complexity

Ian Herbert

National University of Singapore

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Getting on the Same Page

Definition

A *prefix-free machine* is a partial computable function $M : 2^{<\omega} \rightarrow 2^{<\omega}$ such that if $M(\sigma) \downarrow$ then $M(\tau) \uparrow$ for all $\tau \succ \sigma$.

We think of machines as being decoding algorithms.

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The *prefix-free Kolmogorov complexity*, $K(\sigma)$, of a string $\sigma \in 2^{<\omega}$ is the length of the shortest input to the universal prefix-free machine, \mathbb{U} , that produces σ .

Definition/Theorem (Schnorr)

A real A is *Martin-Löf Random* if $\exists b \in \mathbb{N} \forall n \in \mathbb{N} K(A \upharpoonright_n) > n - b$.

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Definition (Chaitin; Solovay)

A real A is *K-trivial* if for all n , $K(A \upharpoonright_n) \leq^+ K(n)$.

Definition (Muchnik)

A real A is *low for K* if for all σ , $K(\sigma) \leq^+ K^A(\sigma)$.

Definition (Zambella)

A real A is *low for MLR* if every Martin-Löf random real, Z , is Martin-Löf random relative to A , i.e. $K^A(Z \upharpoonright_n) >^+ n$

Theorem

- K -trivial \Leftrightarrow Low for K \Leftrightarrow Low for MLR (Nies 2005).
- The K -trivials are closed downward under \leq_T (Nies 2005).
- The K -trivials are closed under effective join (Downey, Hirschfeldt, Nies, Stephan, 2003).
- There are only countably many K -trivials, and they are all Δ_2^0 (Chaitin, 1976).
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One Way to Weaken

Definition

A is Δ_2^0 -bounded K -trivial if for all n , $K(A \upharpoonright_n) \leq^+ K(n) + f(n)$ for all Δ_2^0 orders f .

Definition

A is Δ_2^0 -bounded low for K if for all σ , $K(\sigma) \leq^+ K^A(\sigma) + f(\sigma)$ for all Δ_2^0 orders f .

We use $\mathcal{KT}(\Delta_2^0)$ and $\mathcal{LK}(\Delta_2^0)$ to denote these sets of reals. Why Δ_2^0 ?

Theorem (Baartse, Barmpalias)

There is a Δ_3^0 order g such that $\mathcal{KT}(g)$ is exactly the set of K -trivials.

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There is a Δ_3^0 order g such that $\mathcal{KT}(g)$ is exactly the set of K -trivials.

Theorem

- $\mathcal{LK}(\Delta_2^0) \Rightarrow \mathcal{KT}(\Delta_2^0)$, but the implication does not reverse (H. 2013).
- $\mathcal{LK}(\Delta_2^0)$ contains a perfect set. (H. 2013)
- $\mathcal{LK}(\Delta_2^0)$ is closed downward under \leq_T , but for any real A , there is a $B \in \mathcal{KT}(\Delta_2^0)$ with $A \leq_T B$. (H. 2013)
- $\mathcal{KT}(\Delta_2^0)$ is closed under effective join, but for any real A , there are $B, C \in \mathcal{LK}(\Delta_2^0)$ with $A \leq_T B \oplus C$. (H.)

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Proposition

- *No ML-random is in $\mathcal{LK}(\Delta_2^0)$ or $\mathcal{KT}(\Delta_2^0)$.*
- *If A is Δ_2^0 and in $\mathcal{KT}(\Delta_2^0)$, then A is K -trivial.*
- *$\mathcal{LK}(\Delta_2^0) \Rightarrow$ Low for Effective Dimension. (Hirshfeldt, Weber)*
- *$\mathcal{LK}(\Delta_2^0) \Rightarrow$ Finite Self-Information $\Rightarrow GL_1$ (Hirshfeldt, Weber).*
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‘Strong’ reducibilities like \leq_T , \leq_{tt} , \leq_m have an underlying map:
 $A \leq B$ iff $\exists \Phi : 2^\omega \rightarrow 2^\omega$ with $\Phi(B) = A$.

‘Weak’ reducibilities do not have such an underlying map. The examples we are concerned with all relate to Kolmogorov complexity.

Definition (Downey, Hirschfeldt, LaForte)

$A \leq_K B$ iff for all n , $K(A \upharpoonright_n) \leq^+ K(B \upharpoonright_n)$.

Definition (Nies)

$A \leq_{LK} B$ iff for all σ , $K^B(\sigma) \leq^+ K^A(\sigma)$.

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Since we no longer have an underlying map, uncountably many reals may be reducible to a single real under these reducibilities. A natural questions is:

Question

What are the cardinalities of the lower cones for $\mathcal{KT}(\Delta_2^0)$ in \leq_K and $\mathcal{LK}(\Delta_2^0)$ in \leq_{LK} ?

Definition (Barnali, Vlek)

A real A is *infinitely often K -trivial* if for infinitely many n , $K(A \upharpoonright_n) \leq^+ K(n)$.

Definition (Miller)

A real A is *weakly low for K* if for infinitely many σ , $K(\sigma) \leq^+ K^A(\sigma)$.

Theorem (Barnali, Vlek)

- *Every r.e. set is i.o. K -trivial.*
- *Every \leq_{tt} -degree contains an i.o. K -trivial.*
- *There is a perfect set of i.o. K -trivials.*
- *Every set that is computed by a 1-generic is i.o. K -trivial.*
- *No Martin-Löf random set is i.o. K -trivial.*
- *If A is i.o. K -trivial, then A has a countable lower \leq_K -cone.*

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Theorem (H. with Stephan)

If A is Δ_2^0 -bounded K -trivial, then A is infinitely often K -trivial, and this implication does not reverse.

Corollary

Every real in $\mathcal{KT}(\Delta_2^0)$ has a countable lower \leq_K -cone.

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A is weakly low for K iff A is low for Ω , i.e. $\Omega = \mu(\text{dom}(\mathbb{U}))$ is ML-random relative to A.

Corollary (via Nies, Stephan, Terwijn)

A is 2-random iff A is ML-random and weakly low for K.

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Theorem

- *Weakly Low for K is closed downward under \leq_T .*
- *If A is weakly low for K then it is GL_1 ($A' \equiv_T A \oplus \emptyset'$) (Nies, Stephan, Terwijn).*
- *If A is Δ_2^0 and weakly low for K , then A is low for K (follows from Hirschfeldt, Nies, Stephan).*

And most importantly for us,

Theorem (Barnpalias, Lewis)

A has a countable lower \leq_{LK} -cone if and only if A is weakly low for K .

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Back to Δ_2^0 -Bounded

So do we have that Δ_2^0 -bounded low for K implies weakly low for K , and we can be done?

Unfortunately, no:

Theorem (H.)

Neither of weakly low for K and Δ_2^0 -bounded low for K implies the other.

Corollary

Some Δ_2^0 -bounded low for K reals have countable lower \leq_{LK} -cones, and some have uncountable ones.

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Can we characterize those reals that are both Δ_2^0 -bounded low for K and weakly low for K ?

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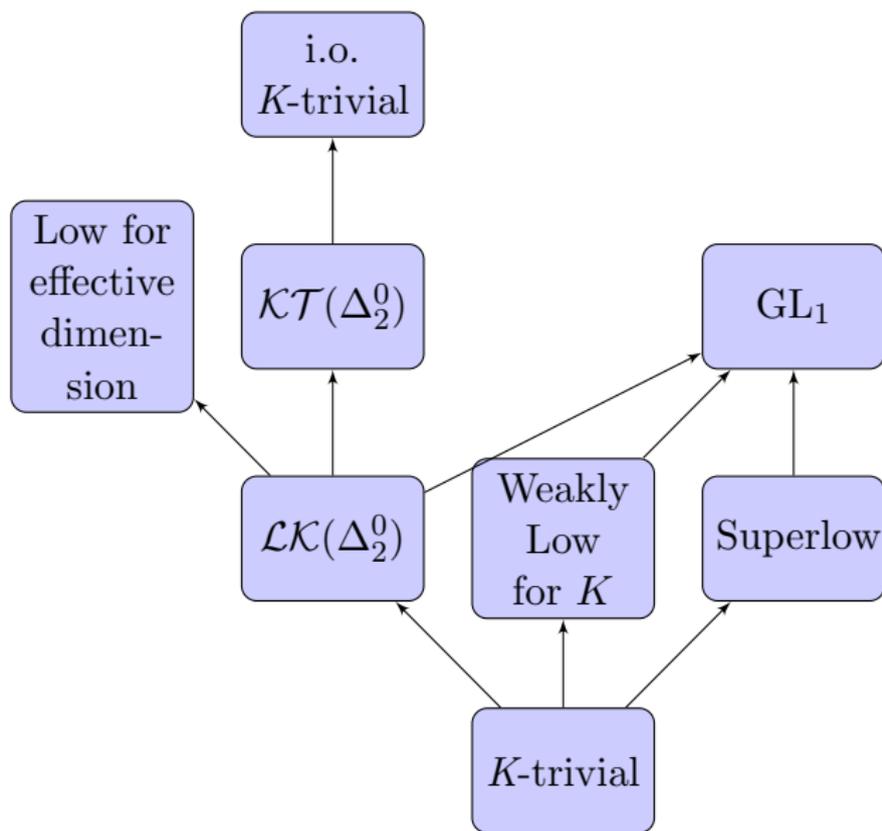
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Question

Every nonrecursive weakly low for K set is of hyperimmune degree (Miller, Nies). What about $\mathcal{LK}(\Delta_2^0)$?

Question

What can we say about the internal structures of $\mathcal{LK}(f)$ and $\mathcal{KT}(g)$ for various f and g under \leq_{LK} and \leq_K ?

Question

What about other lowness notions? C -triviality, lowness for C , etc?

Thanks!